CIS 6930: IoT Security

Lecture 5

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Derived from slides by Adwait Nadkarni, William Enck, Micah Sherr and Patrick McDaniel

Class Notes and Clarifications

- General comments on paper reviews
 - Pay attention to the feedback!
 - The grading metric also looks at your improvements, so not repeating mistakes is key!
 - Small problems in implementation vs larger limitations/drawbacks of the methodology/framework
 - Your comments should match the accept/reject verdict!
- General comment on discussions
 - I expect you to have at least some thoughts on the research topics we discuss!
 - The point is to express opinions and engage.



Class Notes and Clarifications

- General comments on presentations
 - Deeply understand and relay the ideas of the paper that you are presenting
 - If you don't understand something yourself, you won't be able to explain it to us.
 - Understand why the authors are choosing to do things a certain way.
 - Since you are presenting, we'll assume you are the author!
 - Remember that the presentation is a visual medium!

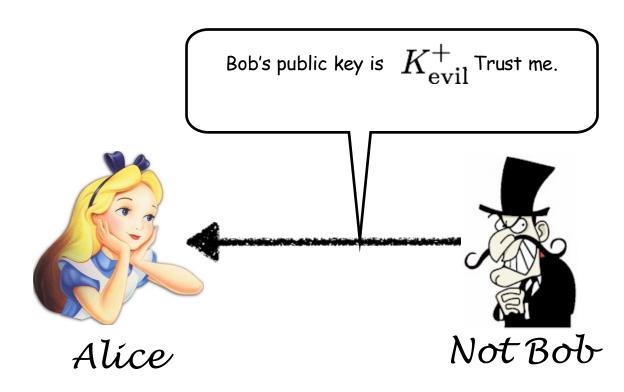


Public Key Crypto (10,000 ft view)

- Separate keys for encryption and decryption
 - Public key: anyone can know this
 - Private key: kept confidential
- Anyone can encrypt a message to you using your public key
- The private key (kept confidential) is required to decrypt the communication
- Alice and Bob no longer have to have *a priori* shared a secret key

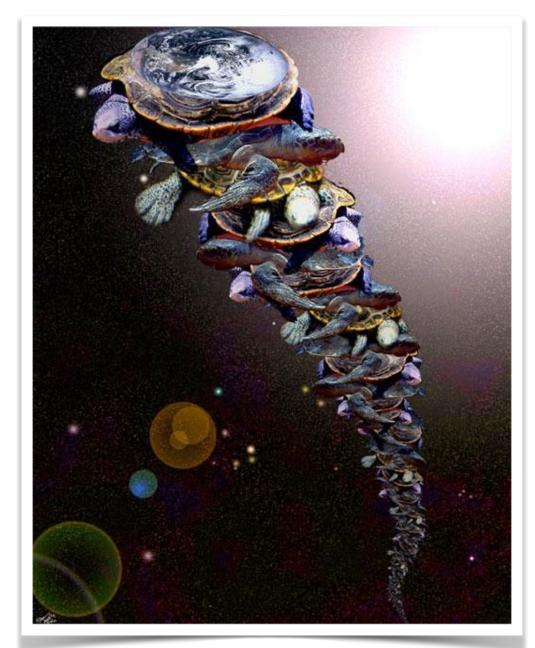
Problem? YES. *How do we know if Alice's key is really Alice's?*

But how do we *verify* we're using the correct public key?



Short answer: We can't.

It's turtles all the way down.



Why not just use a database?

- Every user has his/her own public key and private key.
- Public keys are all published in a database.
- Alice gets Bob's public key from the database
- Alice encrypts the message and sends it to Bob using Bob's public key.
- Bob decrypts it using his private key.
- What's the problem with this approach?

Solving the Turtles Problem

- We need a trust anchor
 - there must be someone with authority
 - requires *a priori* trust
- Solution: form a trust hierarchy
 - "I believe X because..."
 - "Y vouches for X and..."
 - "Z vouches for Y and..."
 - "I implicitly trust **Z**."



Browser Certificate



→ 📴 VeriSign Class	rry Certification Authority 3 Public Primary Certification Authority – G5 lass 3 International Server CA – G3 chase.com	
Certificate Issued by Expires:	hase.com r: VeriSign Class 3 International Server CA – G3 Thursday, August 16, 2012 7:59:59 PM ET ertificate is valid	Î
Details		- 11
Subject Name		- 11
Country	US	
State/Province		
	Jersey City	
-	JPMorgan Chase	
Organizational Unit		
Common Name	www.chase.com	
Issuer Name		- 11
Country	US	
Organization	VeriSign, Inc.	
Organizational Unit	VeriSign Trust Network	
	Terms of use at https://www.verisign.com/rpa (c)10	
Common Name	VeriSign Class 3 International Server CA - G3	
Serial Number Version	61 5C 33 29 65 09 08 60 A4 E6 82 50 00 F6 22 F0 3	
Signature Algorithm Parameters		
Not Valid Before	Tuesday, August 16, 2011 8:00:00 PM ET	
	Thursday, August 16, 2012 7:59:59 PM ET	<u> </u>

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OK

What's a certificate?

- A certificate ...
 - ... makes an association between an identity and a private key
 - ... contains public key information {e,n}
 - ... has a validity period
 - ... is signed by some *certificate authority* (CA)
 - ... identity may have been vetted by a registration authority (RA)
- People trust CA (e.g., Verisign) to vet identity

Why do I trust the certificate?

- A collections of "root" CA certificates (self-signed)
 - ... baked into your browser
 - ... vetted by the browser manufacturer
 - ... <u>supposedly</u> closely guarded
 - trust anchor
- Root certificates used to validate certificate
 - Vouches for certificate's authenticity

Certificate Manager

ertificate Name	tes on file that identify t	nese certifica			
	e dy Group, Inc.		Security D	Device	
	Secure Certification Au	thority	Software	Security Device	ē
	Class 2 CA	litority		ject Token	
The USERTRI			bunch Ob	Jeer roken	
	Solutions Certificate Aut	hority	Software S	Security Device	
	om CA SSL Services (OV			Security Device	
-	RFirst-Hardware	, 		ject Token	
UTN - DA	TACorp SGC			ject Token	
UTN-USE	RFirst-Network Applicati	ons		ject Token	
UTN-USE	RFirst-Client Authenticat	tion and Emai	il Builtin Ob	ject Token	
UTN-USE	RFirst-Object		Builtin Ob	ject Token	
Türkiye Bilim	nsel ve Teknolojik Araştı	rma Kurumu.			
TÜBİTAK	UEKAE Kök Sertifika Hizr	net Sağlayıcı.	Builtin Ob	ject Token	
TÜRKTRUST	Bilgi İletişim ve Bilişim O	üvenliği Hiz.			
TÜRKTRU	ST Elektronik Sertifika H	izmet Sağlay.	Builtin Ob	ject Token	
University of	Pennsylvania				
DSL CA A	uthority		Software S	Security Device	
Unizeto Sp. :	z o.o.				0
Certum C	A		Builtin Ob	ject Token	
ValiCert, Inc.					
RSA Publi	c Root CA v1		Software S	Security Device	
http://ww	vw.valicert.com/		Builtin Ob	ject Token	
http://ww	vw.valicert.com/		Builtin Ob	iect Token	Ŧ

← → C $\land https://www.csc.ncsu.edu$	M	≡

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Your connection is not private

Attackers might be trying to steal your information from **www.csc.ncsu.edu** (for example, passwords, messages, or credit cards). NET::ERR_CERT_COMMON_NAME_INVALID

Automatically report details of possible security incidents to Google. Privacy policy

Advanced

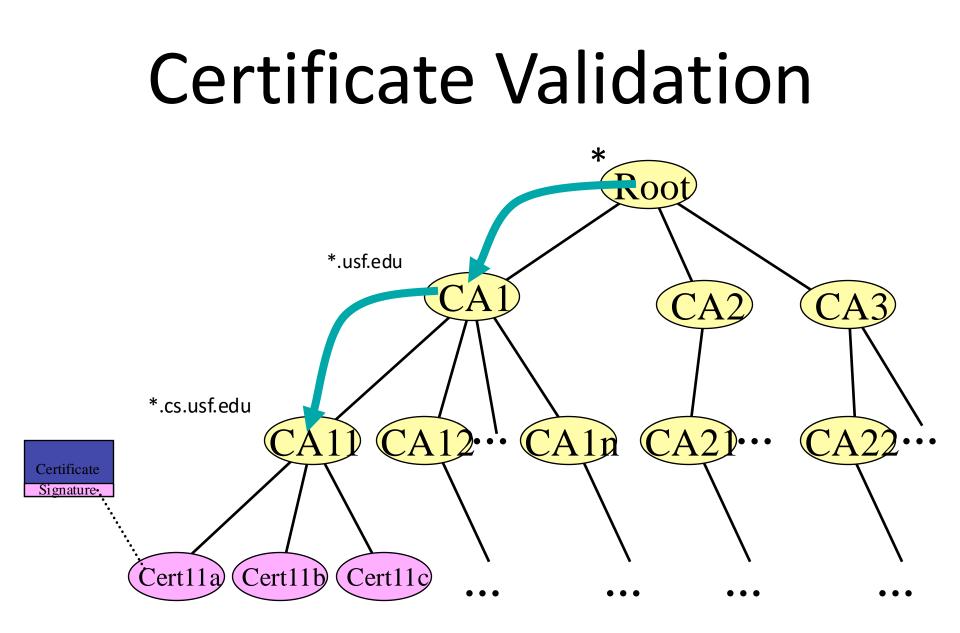
Back to safety

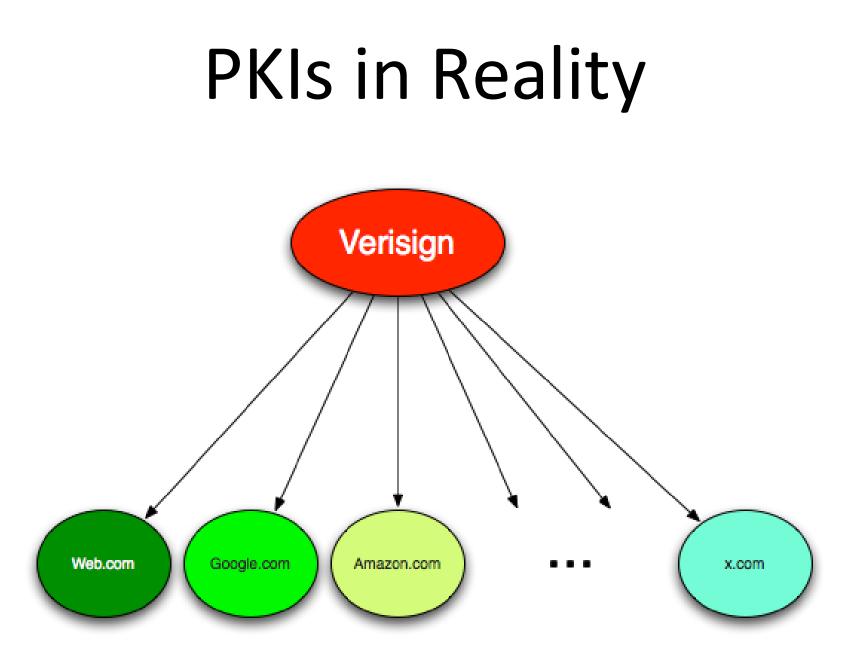
Public Key Infrastructure

- Hierarchy of keys used to authenticate certificates
- Requires a root of trust (i.e., a trust anchor)

What is a PKI?

- Rooted tree of CAs
 - Cascading * Root issuance Any CA can *.usf.edu *.chase.com issue cert CA CAs issue certs CA CA for children *.cs. usf. edu **CA12**... CA2 CA22CA1n CA] Cert11a (Cert11b (Cert11c)



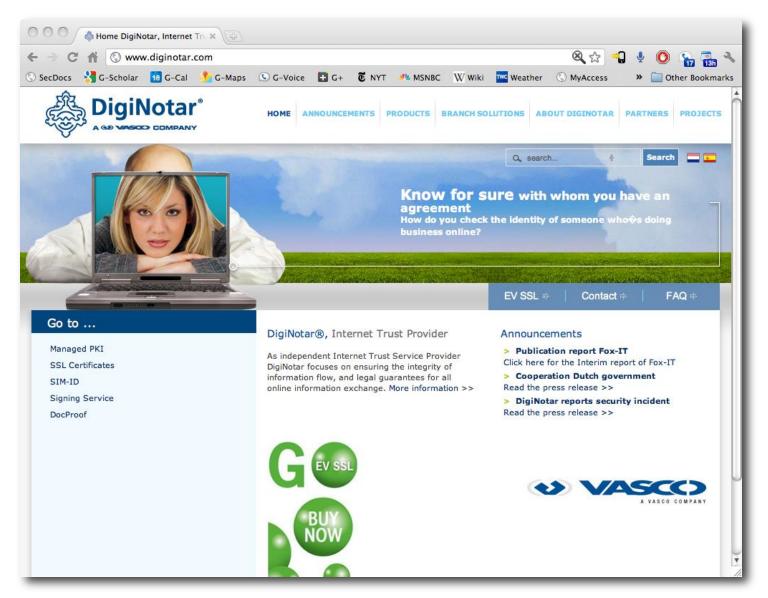


Obtaining a Certificate

- 1.Alice has some identity document A^{ID} and generates a keypair (A⁻, A⁺)
- $\textbf{2.A} \rightarrow \textbf{CA}: \ \{\textbf{A}^{\scriptscriptstyle +}, \textbf{A}^{\scriptscriptstyle \text{ID}}\}, \ \textbf{Sig}(\textbf{A}^{\scriptscriptstyle -}, \{\textbf{A}^{\scriptscriptstyle +}, \textbf{A}^{\scriptscriptstyle \text{ID}}\})$
 - CA verifies signature -- proves Alice has A⁻
 - CA may (and should!) also verify A^{ID} offline
- **3.** CA signs $\{A^+, A^{ID}\}$ with its private key (CA⁻)
 - CA attests to binding between A+ and A^{ID}
- **4.**CA \rightarrow A : {A⁺, A^{ID}}, Sig(CA⁻, {A⁺, A^{ID}})
 - this is the certificate; Alice can freely publish it
 - anyone who knows CA⁺ (and can therefore validate the CA's signature) knows that CA "attested to" {A⁺, A^{ID}}
 - note that CA never learns A⁻

- Any CA may sign any certificate
- Browser weighs all root CAs equally
- Q: Is this problematic?

The DigiNotar Incident



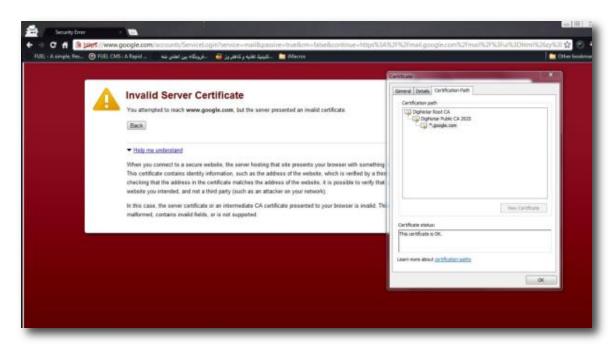
DigiNotar Incident

- DigiNotar is a CA based in the Netherlands that is (well, was) trusted by most OSes and browsers
- July 2011: Issued fake certificate for gmail.com to site in Iran that ran MitM attack...
- ... this fooled most browsers, but...



DigiNotar Incident

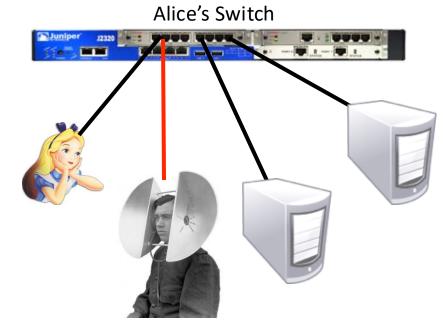
- As added security measure, Google
 Chrome hardcodes
 fingerprint of
 Google's certificate
- Since DigiNotar didn't issue Google's true certificate, this caused an error message in Chrome

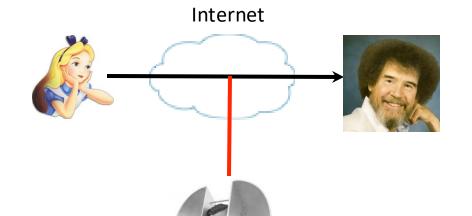


How secure is the verifier?

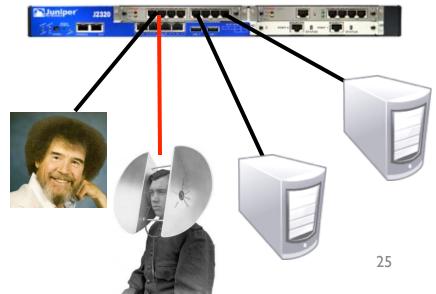
- What happens if attacker is able to insert his public root CA key to the verifier's list of trusted CAs?
- More generally, what are the consequences if the verifier is compromised?
- Q: What's the consequences for IoT devices/apps?

Eavesdropping





Bob's Switch



Why is crypto useful?

		ter and the second seco	tcat_unencrypted.pcapng	
		🗙 🌀 🍳 🦛 🛒		
MacBook-Pro-4	<pre>\$ echo "Securit</pre>	y is Fun" ne	etcat -v localhost	Expression + Apply this filter
localhost [127.0.0.1]	8080 (http-alt)	open		
		open		2 Ack=2 Win=12758 Len=0 TSv
				Seq=2 Ack=22 Win=12758 Len=0 CWR] Seq=0 Win=65535 Len=0
	LOG EM LETTOTOTI	12,101011		
	160 2 127.0.0.1	127.0.0.1	TCP 56 59584 → 8080	0 [ACK] Seq=1 Ack=1 Win=408288 Len=0 TSv
	161 2 127.0.0.1	127.0.0.1		Update] 8080 \rightarrow 59584 [ACK] Seq=1 Ack=1
	162 2 127.0.0.1 163 2 127.0.0.1	127.0.0.1 127.0.0.1		<pre>P [PSH, ACK] Seq=1 Ack=1 Win=408288 Len= P [ACK] Seq=1 Ack=17 Win=408256 Len=0 TS</pre>
	164 2 127.0.0.1	127.0.0.1		36 [SYN] Seq=0 Win=65535 Len=0 MSS=16344
	165 2 127.0.0.1	127.0.0.1		35 [RST, ACK] Seg=1 Ack=1 Win=0 Len=0
	166 2 127.0.0.1	127.0.0.1		36 [SYN] Seq=0 Win=65535 Len=0 MSS=16344
	167 2 127.0.0.1	127.0.0.1		86 [RST, ACK] Seq=1 Ack=1 Win=0 Len=0
	168 2 127.0.0.1	127.0.0.1		36 [SYN] Seq=0 Win=65535 Len=0 MSS=16344
	169 2 127.0.0.1 170 2 127.0.0.1	127.0.0.1 127.0.0.1		87 [RST, ACK] Seq=1 Ack=1 Win=0 Len=0 86 [SYN, ECN, CWR] Seq=0 Win=65535 Len=0
	170 2 127.0.0.1 171 2 127.0.0.1	127.0.0.1		38 [RST, ACK] Seq=1 Ack=1 Win=0 Len=0
	172 2 127.0.0.1	127.0.0.1		36 [SYN] Seq=0 Win=65535 Len=0 MSS=16344
	173 2 127.0.0.1	127.0.0.1		39 [RST, ACK] Seq=1 Ack=1 Win=0 Len=0
	ame 162: 72 bytes on wire (576 hits), 72 hytes ca	aptured (576 bits) on interfa	ace 0
	ll/Loopback	5/6 51(5/, /2 5)(C5 C		
	ternet Protocol Version 4,	Src: 127.0.0.1, Dst: 1	127.0.0.1	
► Tra	ansmission Control Protocol	, Src Port: 59584 (595	584), Dst Port: 8080 (8080),	Seq: 1, Ack: 1, Len: 16
0000	02 00 00 00 45 00 00 44	2a ch 40 00 40 06 00 0	0ED *.@.@	
	7f 00 00 01 7f 00 00 01			
0020 0030	d9 cd 38 c7 80 18 31 d7 6a 50 15 48 6a 50 15 47			
	20 69 73 20 46 75 6e 0a	55 65 75 72 69 74 7	is Fun.	
				1
				WIRESHARK
• 2	1		Packets: 199 · Di	splayed: 199 (100.0%) · Load time: 0:0.3 Profile: Default

• Its just an instant message, right?

Alice uses the Internet for:

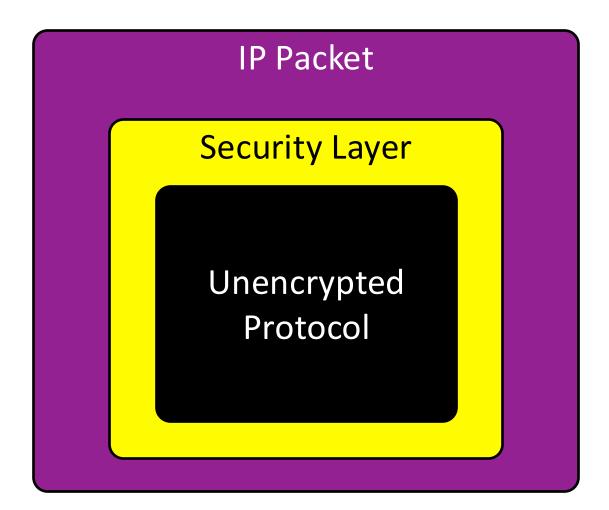
- Email Why is this bad?
- Banking
- Online shopping
- Social networking
- ...

Let's use that crypto stuff

- Let's build some new protocols
 - HTTP \rightarrow SecureHTTP
 - POP \rightarrow POPSecure
 - IMAP \rightarrow Cry
 - SMTP \rightarrow SMTP
 - FTP \rightarrow FTPS
 - Jabber \rightarrow SecJabber
 - Telnet \rightarrow TelCryptNet

Let's build a cryptowrapper standard instead





What properties should this crypto-wrapper have?

- Confidentiality
- Integrity
- Authenticity
 - Server
 - Client
 - Mutual authentication

SSL / TLS

History

- Secure Sockets Layer (SSL) developed by Netscape (remember them?) in 1995
 - Version 1 never released
 - Version 2 incorporated into Netscape Navigator 1.1
 - Microsoft fixes vulnerabilities in SSLv2 and introduces Private Communications Technology (PCT) protocol
 - Netscape overhauls SSLv2, fixing some more security issues, and releases SSLv3
 - IETF takes over and releases Transport Layer Security (TLS), a noninteroperable upgrade to SSLv3
 - current version is TLS version 1.3, <u>RFC 8446</u> (August 2018)

K.I.S.S.

- Application-layer protocol
- Operates over TCP --WHY?



Overview

- Alice (client) initiates conversation with Bob (server)
- Bob sends Alice his certificate
- Alice verifies certificate
- Alice picks a random number S and sends it to Bob, encrypted with Bob's public key
- Both parties derive key material from S
- Client and server exchange encrypted and integrityprotected data

SSLv2 Handshake

ClientHello, Version, Cipher list., RAlice

ServerHello, Ver., Cert., Chosen cipher, RBob

<u>E(K',Data)</u>

E(Bob+,S)



Alice computes master secret k as K=f(S,R_{Alice},R_{Bob})

Hice

Encryption and integrity keys derived from Master Secret

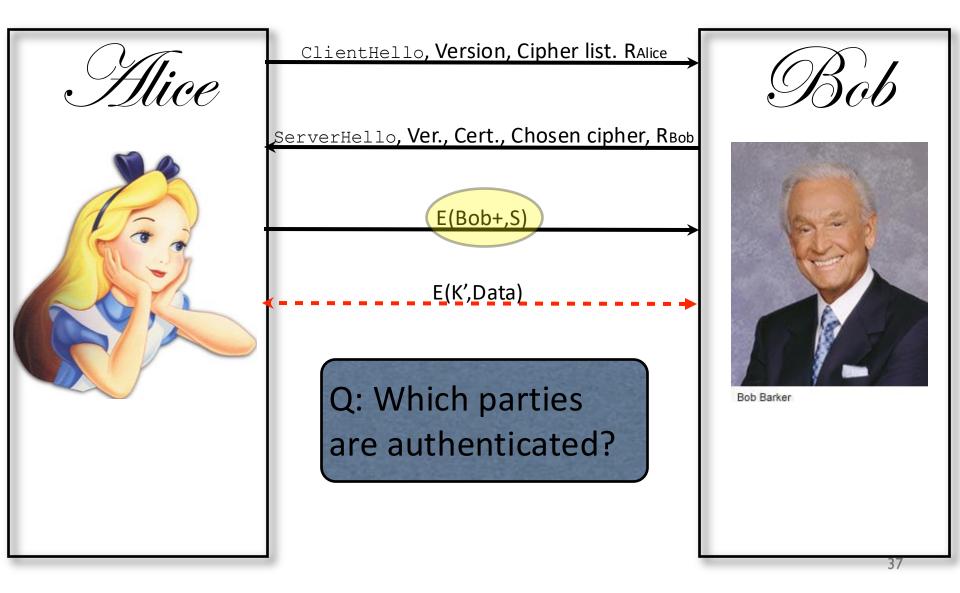
Alice randomly chooses S

Bob computes master secret k as K=f(S,R_{Alice},R_{Bob})

Cryptographic Parameters

- Generated from
 - the master secret K
 - Rc
 - Rs
- Six values to be generated
 - client authentication and encryption keys
 - server authentication and encryption keys
 - client encryption IV
 - server encryption IV
- Generator functions: k_i = g_i(K,Rc,Rs)

Authentication



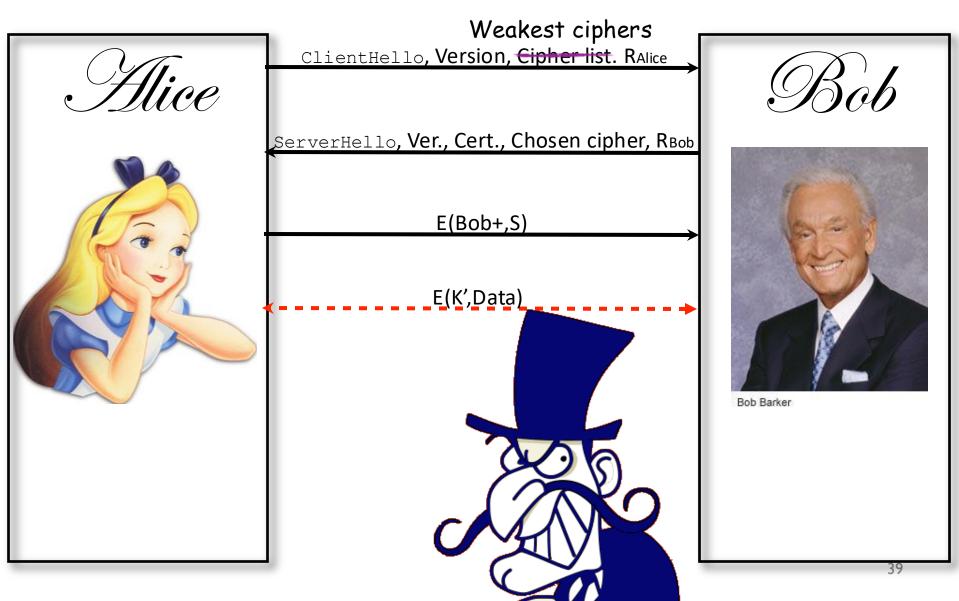
Cipher Suites

- Alice gives Bob a list of supported cipher suites; Bob makes final choice
- Includes encryption algorithms, key length, block mode, and integrity checksum algorithm
- Only 5 supported in TLS1.3, >30 in TLS1.2

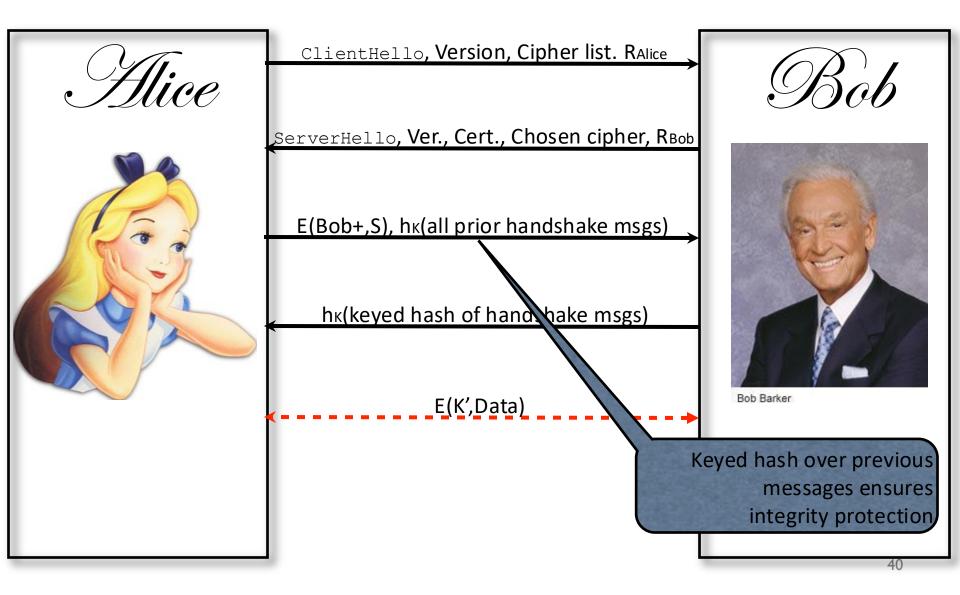
4	
010	openssl ciphers -v
ΓI	LS_AES_256_GCM_SHA384 TLSv1.3 Kx=any Au=any Enc=AESGCM(256) Mac=AEAD
TI	LS_CHACHA20_POLY1305_SHA256 TLSv1.3 Kx=any Au=any Enc=CHACHA20/POLY1305(256) Mac=AEAD
TI	LS_AES_128_GCM_SHA256 TLSv1.3 Kx=any Au=any Enc=AESGCM(128) Mac=AEAD
ЕC	CDHE-ECDSA-AES256-GCM-SHA384 TLSv1.2 Kx=ECDH Au=ECDSA Enc=AESGCM(256) Mac=AEAD
ЕC	CDHE-RSA-AES256-GCM-SHA384 TLSv1.2 Kx=ECDH Au=RSA Enc=AESGCM(256) Mac=AEAD
DH	HE-RSA-AES256-GCM-SHA384 TLSv1.2 Kx=DH Au=RSA Enc=AESGCM(256) Mac=AEAD
ЕC	CDHE-ECDSA-CHACHA20-POLY1305 TLSv1.2 Kx=ECDH Au=ECDSA Enc=CHACHA20/POLY1305(256) Mac=AEAD
ЕC	CDHE-RSA-CHACHA20-POLY1305 TLSv1.2 Kx=ECDH Au=RSA Enc=CHACHA20/POLY1305(256) Mac=AEAD
DH	HE-RSA-CHACHA20-POLY1305 TLSv1.2 Kx=DH Au=RSA Enc=CHACHA20/POLY1305(256) Mac=AEAD

- Key Exchange algos e.g. RSA, DH, ECDH
- Authentication algos e.g., RSA
- Bulk encryption algos e.g., AES
- MAC algos e.g., SHA-256

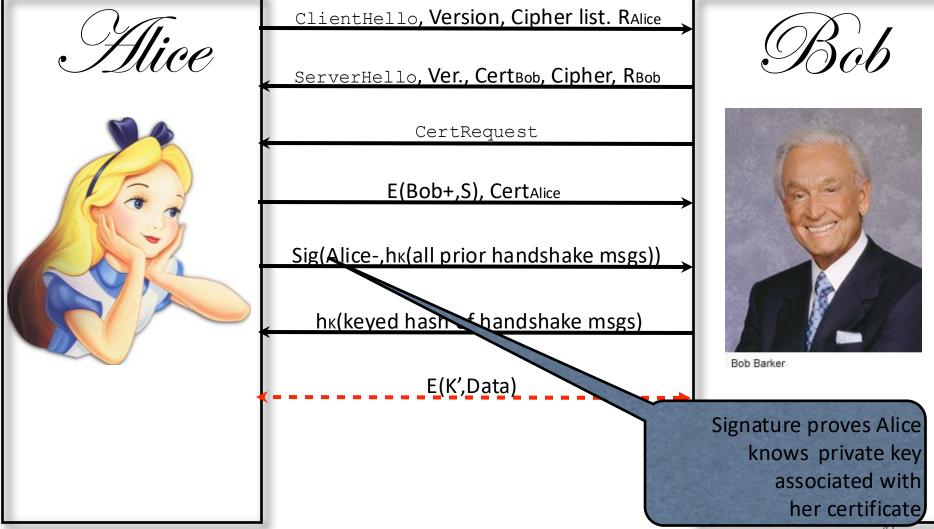
SSLv2 Problems



SSLv3 Fixes

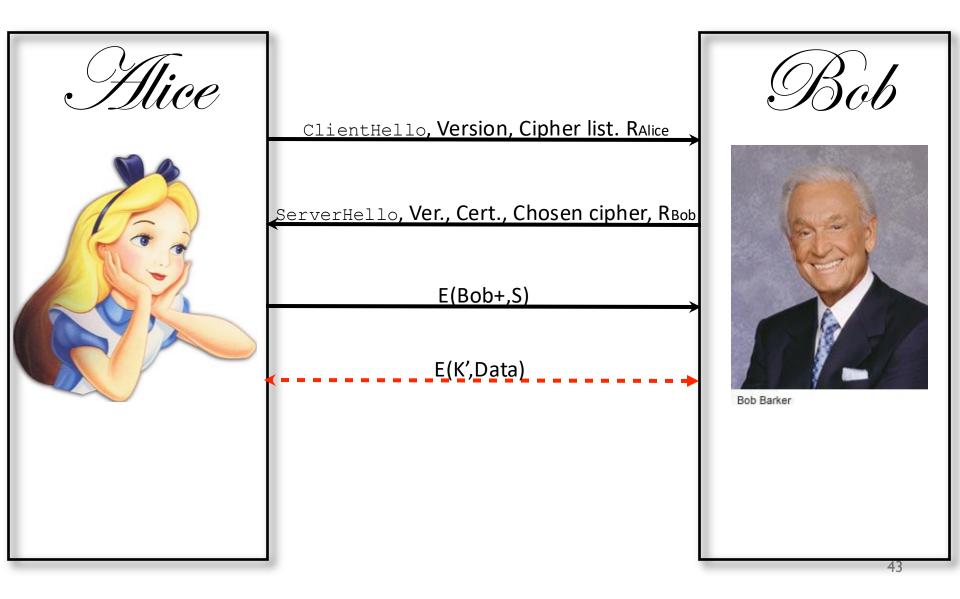


SSL/TLS with Server and Client Authentication

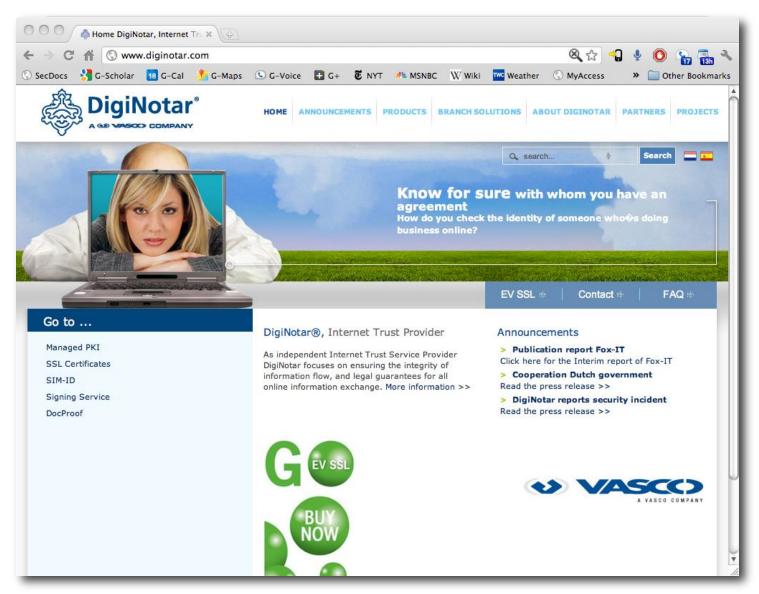


Problems with TLS/SSL

If Bob's cert isn't verified, how do you know you're actually talking to Bob?

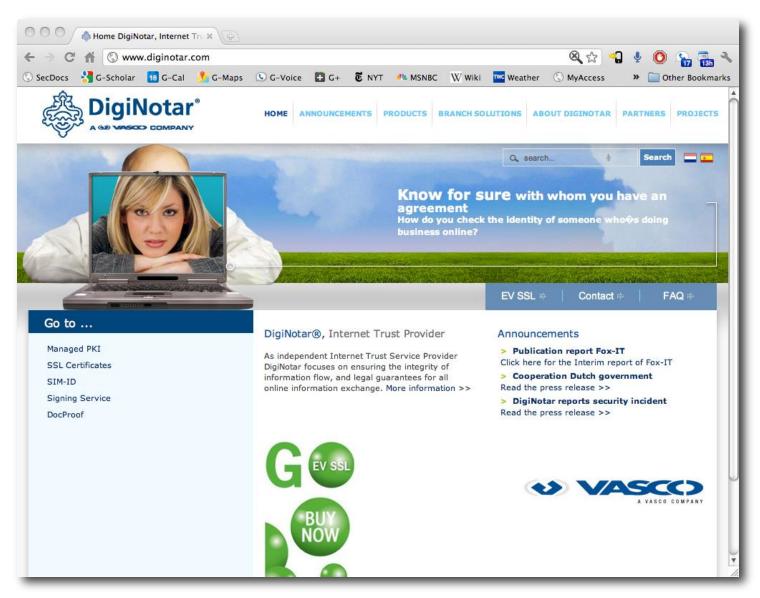


Solution: Use a PKI



- Any CA may sign any certificate
- Browser weighs all root CAs equally
- *Q: Do you recall why this is problematic?*

Recall: The DigiNotar Incident



SSL/TLS in the Real World

Network Stack, revisited

Application
SSL/TLS
Transport
Network
Link
Physical

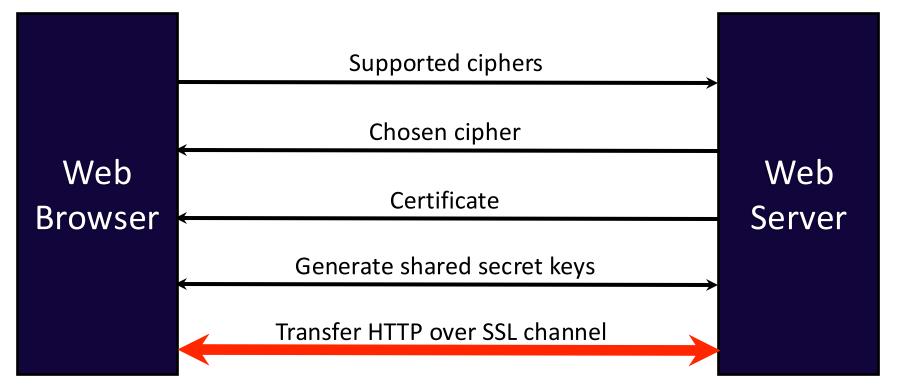
SSL/TLS in the Real World

- All (modern) browsers support TLS 1.2, TLS1.3
 - SSLv3 deprecated in most major browsers
- Client authentication very rare -- WHY?
- Implementations:
 - HTTP (80) \rightarrow HTTPS (443)
 - POP (110) → POP3S (995)
 - IMAP (143) \rightarrow IMAPS (993)
 - SMTP (25) \rightarrow SMTP with SSL (465)
 - FTP (20,21) → FTPS (989,990)
 - Telnet (23) \rightarrow Telnets (992)

SSL/TLS and the Web

- HTTPS: Tunnel HTTP over SSL/TLS
- Add golden lock symbol





The verifier matters

- SSL is an *application layer protocol*
 - Software developers must use it correctly
- Pre-Smartphone World
 - Small set of applications that use SSL (E.g., Web Browser)
 - Lots of attention to those apps

- Smartphone World
 - Possibly millions of applications that use SSL
 - Many apps do not verify certificates correctly – Implications?
 - Developers change default configuration – WHY?

SSL Verification in Apps

- Even popular apps are vulnerable to incorrect SSL use
 - Banking
 - Document storage
 - Social Networks (Facebook, before Firesheep)
 -and IoT apps
 - •
 - Common mistakes: Generally, in HTTPS use.
 - **1.** Not using SSL
 - 2. Mixed SSL use
 - **3.** Accepting all certificates
 - **4.** Accepting all hostnames (i.e., regardless of the CN)
 - **5.** Trusting all CAs

Not using SSL

- What happens when you don't use SSL? E.g., http://www.mybank.com/loggedin?sessionid=11
 - If I can *guess, infer,* or *steal* the session ID, game over
- Are there any use cases where not using SSL would be okay?
 - It depends. However, unless confidentiality and authenticity are *never* going to be important to the app, use SSL!

Lesson I: Always use SSL (i.e., mostly HTTPS)

Mixed SSL use



- Mixed use of HTTP and HTTPS on the same site.
- Use case 1: Login page is not HTTPS, but the login form is submitted to a HTTPS page.
 - MiTM can replace HTTPS links with HTTP (i.e., SSL Stripping)
- Use case 2: Login page is HTTPS, but the <u>rest of the website</u> <u>may be HTTP</u>
 - Unencrypted cookies/session IDs! (e.g., Firesheep)

Lesson 2: Use HTTPS throughout

Certificate Validation

Apps can override the *TrustManager* interface

```
69
                                                                                    https://stackoverfl
       SSLContext sslContext = SSLContext.getInstance("SSL");
                                                                                    ow.com/questions
       // set up a TrustManager that trusts everything
                                                                                    /2703161/how-to-
       sslContext.init(null, new TrustManager[] { new X509TrustManager() {
                   public X509Certificate[] getAcceptedIssuers() {
                                                                                    ignore-ssl-
                          System.out.println("getAcceptedIssuers ========");
                                                                                    certificate-errors-
                          return null;
                   }
                                                                                    in-apache-
                                                                                    httpclient-4-0
                   public void checkServerTrusted(X509Certificate[] certs,
                                  String authType) {
                          System.out.println("checkServerTrusted ========");
       } }, new SecureRandom());
```

What is wrong with this example? It accepts all server certificates!

Lesson 3: Always validate the server's certificate

Using self-signed certificates

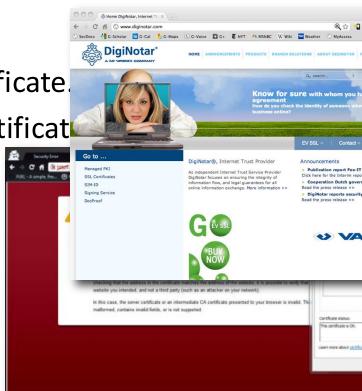
The right way: Certificate Pinning

• i.e., hardcode your self-signed certificate.



Using self-signed certificates

- The right way: Certificate Pinning
 - i.e., hardcode your self-signed certificate.
 - Allows secure use of self-signed certification
- Variation:
 - Pinning own CA certificate
 - Gives you more flexibility.
- How to change the certificate?
 - App updates!
- Don't have to trust 100s of Root CAs!



Lesson 4: Certificate pinning, if done correctly, is more secure than *default SSL use*.

Hostname Verification

- Back to basics: What does a certificate provide?
 - Binding between a *public key* and *identity*

HostnameVerifier hostnameVerifier = org.apache.http.conn.ssl.SSLSocketFactory.ALLOW_ALL_HOSTNAME_VERIFIER;

```
DefaultHttpClient client = new DefaultHttpClient();
```

```
SchemeRegistry registry = new SchemeRegistry();
SSLSocketFactory socketFactory = SSLSocketFactory.getSocketFactory();
socketFactory.setHostnameVerifier((X509HostnameVerifier) hostnameVerifier);
registry.register(new Scheme("https", socketFactory, 443));
SingleClientConnManager mgr = new SingleClientConnManager(client.getParams(), registry);
DefaultHttpClient httpClient = new DefaultHttpClient(mgr, client.getParams());
```

// Set verifier
HttpsURLConnection.setDefaultHostnameVerifier(hostnameVerifier);

https://stackoverflow.com/questions/2012497/accepting-a-certificate-for-https-on-android?lq=1

- Any certificate issued by any trusted CA will be accepted!
 - i.e., HostName= google.com, but cert has CN=foogle.com?

Lesson 5: Never override the HostNameVerifier